

APR - 18/TE/Insem. - 141

T.E. (Computer Engineering)

SYSTEMS PROGRAMMING AND OPERATING SYSTEM

(2015 Course) (Semester - II) (310251)

Time : 1 Hour]

[Max. Marks : 30

Instructions to the candidates:

- 1) Answer Q1 or Q2, Q3 or Q4, Q5 or Q6.
- 2) Neat diagrams must be drawn wherever necessary.
- 3) Figures to the right indicate full marks.
- 4) Assume suitable data, if necessary.

Q1) a) Differentiate between literal and immediate operand. How assembler handles them? Give examples. [6]

b) Define Assembler Directive. Explain ORIGIN, EQU & LTORG with example. [4]

OR

Q2) a) Consider following Assembly code and show output of pass-1 of two pass assembler with entries in MOT, POT, ST, LT and BT. [5]

PROG START 50

USING PROG+2, 15

L 1, FIVE

A 1, = F '2'

LTORG

ST 1, RES

FIVE DC F '4'

RES DS IF

END

b) Discuss with example what is forward reference? How is it handled in single pass assembler? [5]

P.T.O.

- Q3) a) Define Macro. What are the advantages of macro facility? How are macros different from functions? [6]
- b) What is the need of DLL? Differentiate between How Dynamic and static linking? [4]

OR

- Q4) a) What are the types of loaders? Explain compile and go loader scheme with advantages and disadvantages using suitable diagram. [6]
- b) Comment on the statement "Programs with macros require more space and less time at runtime than programs with functions". [2]
- c) Discuss four different functions of loader. [2]

- Q5) a) What is interpreter? Explain various components of interpreter? [4]
- b) Consider input "d = a + b * 2;" and show the output of each phase of compiler with suitable diagram? [6]

OR

- Q6) a) Write regular expressions to recognize following: [4]
- Signed and unsigned integer numbers.
 - Identifiers.
 - Few Keyword in "C" program.
 - Relational Operators.
- b) What is LEX? Explain working of LEX with suitable diagram? [4]
- c) What is Syntax Error? Give suitable example? [2]



SYSTEM PROGRAMMING & OPERATING SYSTEM

INSEM MODEL ANSWER - 2018

08/03/2018

Q.1 a) Differentiate between literal & Immediate operand. How assembler handle them? Give e.g. [6 Marks]



[4-marks]

Sr No	Literal	Sr No	Immediate Operand.
1.	Assembler generates specified value as a constant at some other mem. location.	1.	The operand value is assembled as part of the m/c instruction.
2.	Example: MOVER B, '=5'	2.	Example: MOVE A, #5
3.	Literal use '=' symbol	3.	Immediate Operand use '#' sym
4.	There is mem. reference	4.	There is no mem reference
5.	LITORG ^(AD) instr ⁿ is needed	5.	LITORG Assembler directive is not needed

- Assembler recognise the immediate operand in the instruction by scanning # symbol with constant value.

e.g. MOVER AREG, #5, Y=5

[2-marks]

- Assembler determines the literal in the instrⁿ by scanning '=' symbol with constant value.

e.g. MOVER AREG, '=5'

and

- LITORG instrⁿ is used to recognise the literal from literal table & it makes entry in pool table as total nos of literal.

- Literal value can't be changed during exe.

b) Define Assembler Directive. Explain ORIGIN, EQU & LTORG with example? [04-Marks]

→ Assembler Directives are the instr^s that gives the direction to the assembler, which operation is to be performed. [1 Marks]

1) ORIGIN :-

- This directive gives the direction to assembler to set

Origin location for particular prog/sub-prog. [1 marks]

- Syntax :-

ORIGIN <addr. specification>

where addr. spec^s is an operand, constant.

- This directive sets the address of LC to address given in instr^s.

- e.g. ORIGIN 200

2) LTORG :- [1 marks]

- LTORG instr^s/Directives allows programmer to specify where literal should be placed.

- If LTORG instr^s is not used/present, literal are placed after END directive.

e.g. ^{LC}205 MOVER A, '=2'

206 MOVER B, '=4'

LTORG

LT		PT	
Symb	Addr	0	2
'1'	207	1	-
'2'	208	2	-

Pool table.

- When LTORG instr^s occurred then total nos literal of literal table count & stored entry in pool table. it also assign LC to literal.

3) EQU :- [1-marks]

- EQU directive allow to set equate address to variable/label.

- Syntax is :-

<symbol> EQU <address specification>

where addr. spec^s can be operand or constant.

- The EQU simply associates symbol with the addr. spec^s

e.g.

BACK EQU LOOP

The symbol back is set to address of loop.

Q. 2 a) Consider the following Assembly code and show o/p.

Pass-1 of two pass assembler with entries in MOT, POT, ST, LT & BT. [5-Marks]

PROG START 50

USING PROG+2, 15

LL, FIVE

AL, =F'2'

LTORG

ST 1, RES

FIVE DC F'4'

RES DS IF

END

→

		ST	LT	PT	
	Symbol	Address	Literal address	0	1
50	USING PROG+2, 15				
51	LL, FIVE	55	=F'2'	53	
52	AL, =F'2'	56			
	LTORG				

54 ST 1, RES

55 FIVE DC F'4'

56 RES DS IF

[02-marks]

[02-marks]

MOT

Mnemonics opcode

class

Opcode

length

PROG START 50

AD

01

USING PROG+2, 15

IS

-

LL, FIVE

RG

01

AL, =F'2'

RG

02

LTORG

AD

05

ST 1, RES

IS

01

FIVE DC F'4'

DL

02

RES DS IF

DL

01

END

AD

02

[01-marks]

Q. 2 Discuss with example what is forward References?

How is it handled in Single Pass Assembler? [05-marks]

→ forward references means that variable define in the program before the declaration then we can say it is forward reference.

e.g. START 100

MOVER AREG, X

=====
=====
=====

X DC '1'

A forward reference
to variable X.

[02-marks]

- As assembler can't generate mlc code for assembly instr^s with forward references.

- It can be generated after address of variable used in instr^s is known

- Symbol table is used to store addresses of variables.

These addresses can be used during gener^g of mlc code.

How to Handle forward References.

e.g. START 100

100 MOVER AREG, X

101 ADD BREG, ABC

102 COMP BREG, XY2

103 STOP

104 X DC '2'

105 ABC DC '1'

106 XY2 DC '5'

107 END

Forward
References

Symbol	address
X	104
ABC	105
XY2	106

[03-marks]

- In above e.g. you can see X, ABC & XY2 are forward references variable, which are stored in symbol table &

- address are assigned to it as variable is declared

- If variable is used in instr^s & value is not available, then such condⁿ is called as Backpatching.

Q-3 b) what is need of DLL. Difference betⁿ static & Dynamic Linking.

→ - It stands for Dynamic Link Library [04-Marks]

- Dynamic Link Library is microsoft implementation of:

Shared library in window. [02-Marks]

- The file format for DLL are same as windows EXE. file.

- A DLL can contain code, data & Resources.

- Need of DLL:

- To share code in program, DLL is required.

- shared code is placed in single separate file.

- DLL code may be shared but data is private.

- DLL allow interprocess commⁿ thro. shared memory.



DLL calling a funⁿ in an application.

- DLL may control several clients, when event occurs from DLL.

* Compare static & Dynamic Linking [02-Marks]

Sr no	Static Linking	Sr NO	Dynamic Linking.
1.	Static linker takes object files produced by compiler including lib. fun ⁿ & produces an executable file.	1.	Dynamic linking defers much of linking process until a prog. starts running.
2.	static linking can be fixed.	2.	Dynamic linking can be variable.
3.	No relocation during run time.	3.	Perform reloc ⁿ during run time
4.	Linking is done at compile time	4.	Linking is done at run time.
5.	there is wastage of memory	5.	There is no wastage of memory
6.	It does not support sharing of libraries/memery	6.	Dynamic linking provides automatic sharing of code.

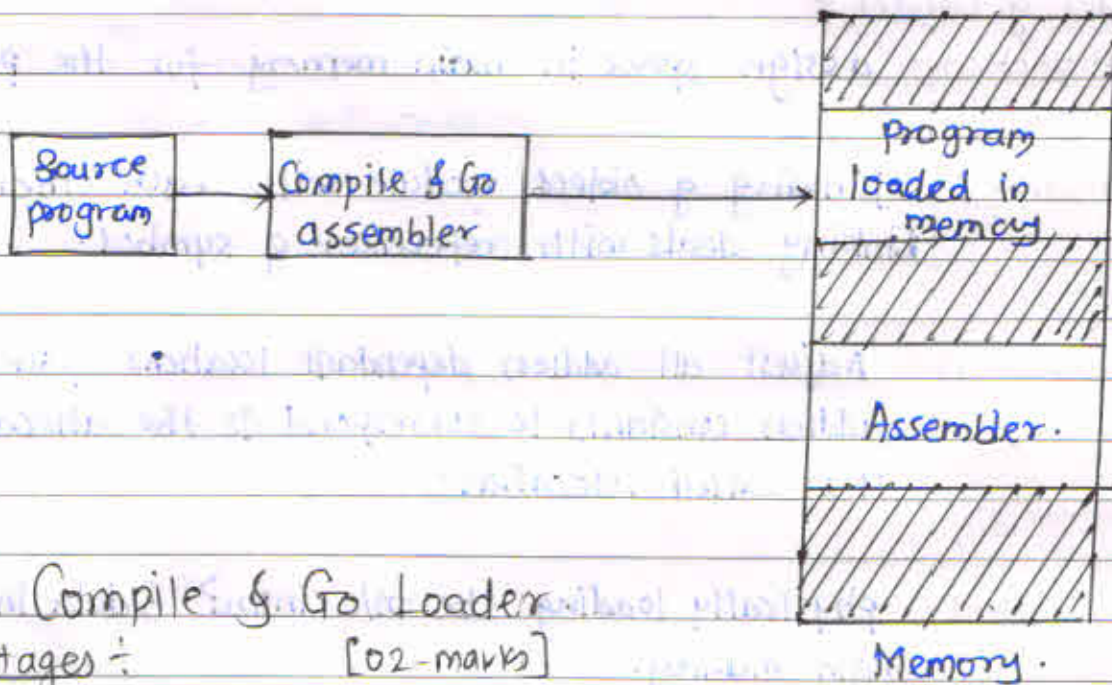
Q. 4 a) What are the types of Loaders? Explain Compile & Go loader schemes with advantages & Disadvantages. [6-Marks]

→ Types of Loaders:-

- 1) Compile & Go loader
- 2) General loader
- 3) Absolute loader
- 4) Subroutine linkage
- 5) Relocating loader
- 6) Direct linking loader.

- Compile & Go loader:- [02-marks]

- In this, Assembler is loaded in one part of memory and it places the assembled prog. (mic instr) directly into their assigned memory location.
- After loading process completes, the assembler transfers the control to starting instr of the loaded prog.



Compile & Go Loader

- Advantages :- [02-marks]

- 1) easy to implement
- 2) It is very simple scheme.
- 3) Assembler available in memory.

- Disadvantages :-

- 1) A portion of memory is wasted, as it is occupied by assembler.
- 2) It is difficult to handle multiple segments.
- 3) It is required to retranslate user prog. every time.
- 4) Difficult to develop modular program.

Q.4 b) Comment on the statement "Program with macro requires more space and less time at runtime than programs with functions" [02-marks]

→ As macro declared as once and used multiple time in prog. by calling with its name.

- macro requires more space because when macro is called in program then macro is expanded so seq. of instr^s can be increased (more space required)

- macro can be called using its name in program and macro is also present in same program so that less time required to call macro rather than function. [02-marks]

Q.4 c) Discuss 4 different function of Loader. [02-marks]

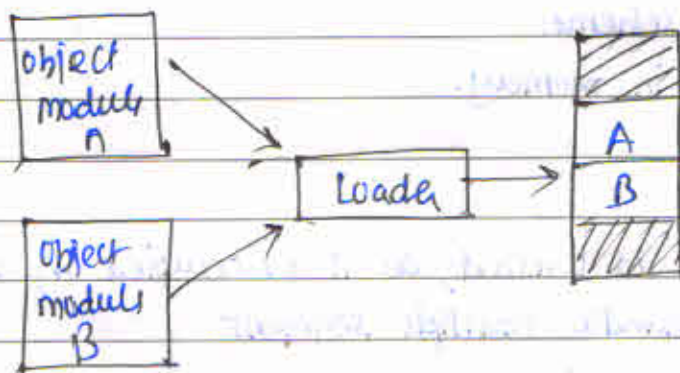
→ Fun^s of Loader :-

1) Allocation - assign space in main memory for the prog.

2) Linking - Linking of object module with each other. Linking deals with references of symbol.

3) Relocation :- Adjust all address dependant locations, such as address constants to correspond to the allocated space. It is called relocation.

4) Loading = physically loading the m/c instr^s & data into main memory. [02-marks]



General Loader Scheme.

Q.5. a) what is Interpreter? Explain Various Component of Interpreter?

→ "Interpreter is a program/tool that which reads the source code one instruction or line at a time, convert this line into machine code or some intermediate form & execute it."

* Component of Interpreter: [01-mark]

1) Assembler 3) Linker 5) Debugger

2) Loader 4) preprocessor [03-marks]

1) Assembler: is a program which converts Assembly lang. into machine level language (code)

2) Loader - it loads the prog. into main memory for the execution.

3) Linker - it links multiple module/ prog. together for exe.

4) Debugger: is used to debug the program line by line.

Q.5 b) Consider input "d = a + b * 2;" and show the o/p of

each phase of compiler with suitable diagram? [06-marks]

→ 1) Lexical Analysis

d = a + b * 2 ;

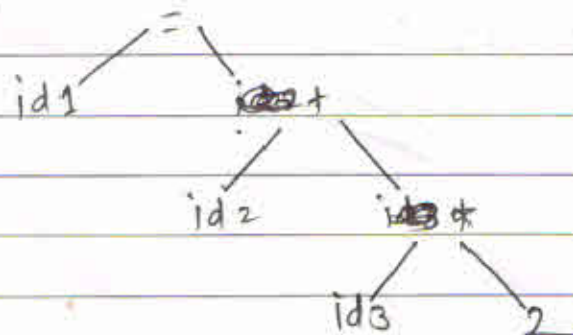
[01-mark]

id1 = id2 + id3 * 2 ;

it scan the from left to right & generates token

2) Syntax Analysis

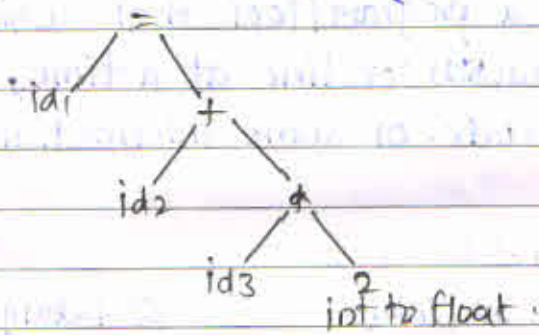
[01-mark]



it checks the syntax of prog.

3) Semantic Analysis:-

[01-mark]



It perform type checking & conversion.

4) Intermediate Code Generation:-

[01-mark]

t1 = Intoreal (2)

t2 = id3 * t1

t3 = id2 + t2

id1 = t3

5) Code Optimizer :-

[01-mark]

t1 = id3 * 2.0

id1 = id2 + t1

6) Code Generator :-

[01-mark]

MOVF R1, id3

MULF R1, #2.0

MOVF R2, id2

ADDF R1, R2

MOVF id1, R1

//

Q6. c) What is Syntax Error? Give Suitable Example. [02-marks]

→ "Syntax error is used checks the Syntan of program whether it is properly defined or not" [02-marks]

- This oper^s is done in 2nd phase of compile. i.e. Syntax analysis

- It is also called as syntux tree.

- It gives various types of error.

e.g. lvalue, rvalue not declared,

Semicolon not defined,

e.g.

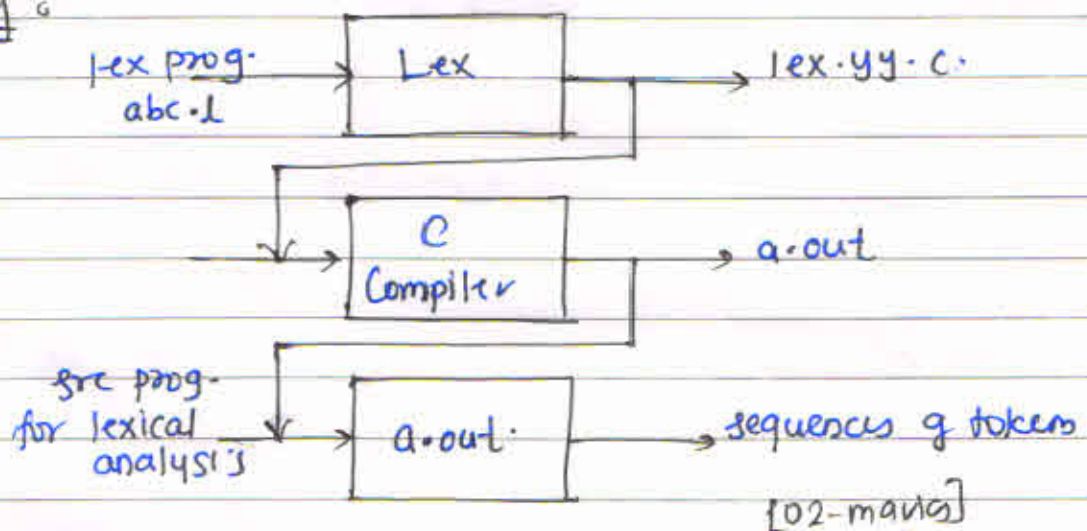
$x = 4 + 2 * 30$

semi-colon missing error.

Q6. b) What is LEX? Explain working of LEX with Suitable example? [04-marks]

→ "LEX is lexical Analyzer generator tool, which is used to scan the program from left to right & seperate out the token to generate parse tree." [01-marks]

Working



- Lex is slow tool, which takes i/p as regular expression together with action to be taken on recognising each of these pattern.

- The o/p of lex is C program for lexical analysis.

- .l prog is to be passed to Lex. it compile lex file & generate C files. then C compiler compiles C file & generate exe files.

- & finally generates token using a.out cmd [01-marks]

Q 6. b) Syntax of Lex

How to run LEX code

→ declaration

> Lex abc.l

%.y.

> gcc cc.y.c.

translation Rule.

> a-out <l

User function

Q 6 a) write Regular expression to recognise following.

1) Signed & Unsigned Integer number

2) Identifier

3) Few keyword in 'C' prog

4) Relational Operator.

[04-marks]

→ 1) The set of strings accepted by finite automata is known as regular lang.

RE = 1, 2, 3, 4.

RE = -1, -2, -3

[01-marks]

2) RE = (id1 + id2 + id3)*

[01-marks]

3) RE = (for + while + do)

[01-marks]

4) RE = (= + > + < + >= + <=)*

[01-marks]

Note: Question No. 6 a), you can solve as per concept of Theory of Computation.

